

Representing And Measuring Networks

Rogayeh Dastranj Tabrizi

email: rda18@sfu.ca

Office: WMC 3607

Office Hours: Thursdays 12pm-2pm

Department of Economics
Simon Fraser University

Thanks to Matthew Jackson for access to his teaching resources.

10. Februar 2015

Simplifying The Complexity

- Global patterns of networks:
 - Degree distributions
 - Path lengths
- Segregation Patterns: node types and homophily
- Local Patterns
 - Clustering
 - Support
- Positions in networks
 - Neighbourhoods
 - Centrality, influence, ...

Representing Networks

- $N = \{1, 2, \dots, n\}$ is the set of **nodes**, or **vertices**, **players**, **agents**.
- Connection between nodes is called **links**, or **edges**, **ties**:
 - 1 They may have **intensity** (**weighted network**):
 - How many hours do people spend together per week?
 - How much of one country's GDP is traded with another?
 - 2 They may just be **0 or 1** (**unweighted network**):
 - Have two researchers written an article together?
 - Are two people "friends" on some social platform?
 - 3 They may be **Directed** or **Undirected**:
 - coauthors, friends, ..., relatives, spouses, ..., are mutual relationships
 - link from one web page to another, citations, following on social media..., one way

Graphs and Networks

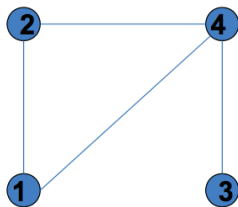
- $N = \{1, 2, \dots, n\}$ is the set of **nodes**,
- $g_{n \times n}$ is a real-valued $n \times n$ matrix, where g_{ij} represent the relationship between i and j .
- In an **unweighted** network:

$$g_{ij} = \begin{cases} 1 & \text{if } ij \in g \\ 0 & \text{otherwise} \end{cases}$$

- Notation: $ij \in g$ indicates a link between i and j .
- Self-links or loops often do not have any real consequences or meaning. Unless otherwise is indicated, **assume** $g_{ii} = 0$.

Unweighted Undirected Networks

$$g = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

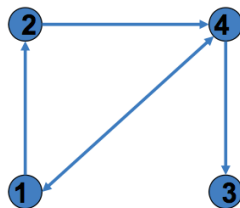


Or list the links:

- $g = \{\{1, 2\}, \{1, 4\}, \{2, 4\}, \{3, 4\}\}$
- $g = \{12, 14, 24, 34\}$ or $g = \{21, 41, 42, 43\}$

Unweighted Directed Networks

$$g = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

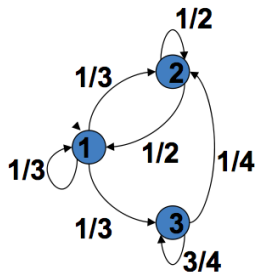


Notice that the **order of nodes** matter in a directed network

- $g = \{12, 14, 41, 24, 34\}$

Weighted Directed Networks

$$g = \begin{pmatrix} 1/3 & 1/3 & 1/3 \\ 1/2 & 1/2 & 0 \\ 0 & 1/4 & 3/4 \end{pmatrix}$$



Graphs and Networks

- A network is **directed** if it is possible that $g_{ij} \neq g_{ji}$.
- A network is **undirected** if $g_{ij} = g_{ji}$ for all nodes i, j .
- $g' \subset g$ indicates that:

$$g' \subset g \Leftrightarrow \{ij : ij \in g'\} \subset \{ij : ij \in g\}$$

- $g + ij$ indicates that a new network that is obtained by **adding** link ij from network g .
- $g - ij$ indicates that a new network that is obtained by **deleting** link ij from network g .

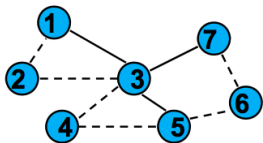
Basic Definitions

- **Walk from i_1 to i_K** : A sequence of links $\{i_1i_2, i_2i_3, \dots, i_{K-1}i_K\}$ such that $i_{k-1}i_k \in g$ for all k in this walk.
- Often it is convenient simply to represent a **walk** as the corresponding sequence of nodes (i_1, i_2, \dots, i_K) such that $i_{k-1}i_k \in g$ for each k .
- **Path from i_1 to i_K** : is a **walk** where all the nodes are distinct.

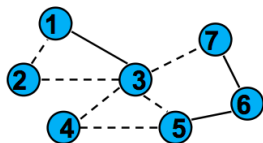
In other words, a **walk** may come back to a given node more than once, whereas a **path** is a walk that never hits the same node twice.

- A **cycle** is a **walk** where $i_i = i_k$.
- **Geodesic**: a **shortest path** between two nodes.

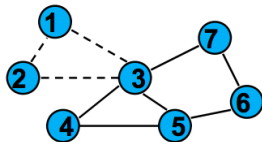
Walks, Paths, Cycles



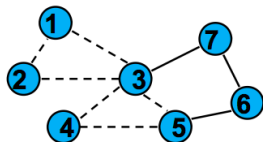
Path (and a walk) from 1 to 7:
1, 2, 3, 4, 5, 6, 7



Walk from 1 to 7 that is not a path:
1, 2, 3, 4, 5, 3, 7



Simple Cycle (and a walk) from 1 to 1:
1, 2, 3, 1

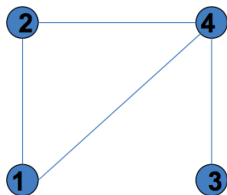


Cycle (and a walk) from 1 to 1:
1, 2, 3, 4, 5, 3, 1

Counting Walks

$$g = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$g^2 = \begin{pmatrix} 2 & 1 & 1 & 1 \\ 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 3 \end{pmatrix}$$

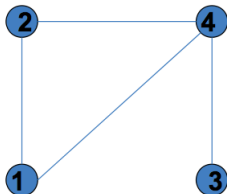


number of walks of length 2 from i to j

Counting Walks

$$g = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$g^3 = \begin{pmatrix} 2 & 3 & 1 & 4 \\ 3 & 2 & 1 & 4 \\ 1 & 1 & 0 & 3 \\ 4 & 4 & 3 & 2 \end{pmatrix}$$

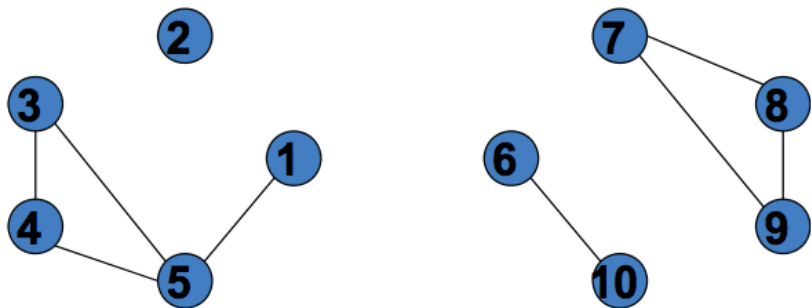


number of walks of length 3 from i to j

Components

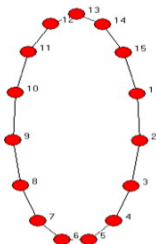
- In many applications it is important to track which nodes can reach which other nodes through a path.
 - Contagion, Learning, Diffusion of various behaviours, etc.
- Network (N, g) is **Connected** if every two nodes are connected by some path.
- **Component**: Maximal connected subgraph, i.e. (N', g') is a component of (N, g) such that:
 - $N' \neq \emptyset, N' \subset N, g' \subset g,$
 - (N', g') is connected,
 - if $i \in N',$ and $ij \in g,$ then $j \in N'$ and $ij \in g'.$
- A link ij is a **bridge** in a network g if $g - ij$ has more components than $g.$

Components

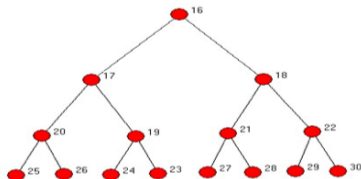


Diameter

- **Diameter**: largest geodesic in the network. If the network is unconnected, then the largest geodesic of the largest component.
- Another measure is **average path length**, which is less prone to outliers.



Diameter is either $n/2$ or $(n - 1)/2$.



Diameter is on the order of $2 \log_2(n + 1)$.

Trees and Stars

There are a few particular network structures that are commonly referred to:

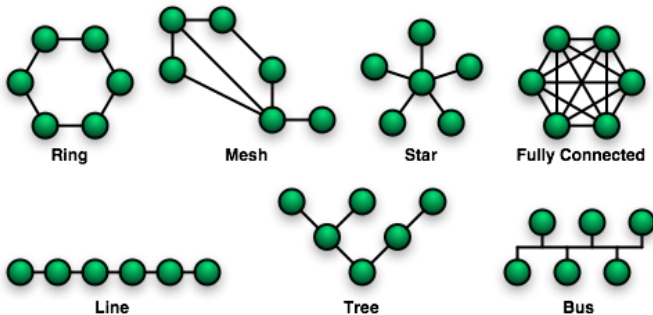
- **Tree:** A connected network that has **no cycles**.
- **Forest:** A network such that each **component** is a tree. Any network that has no cycle is a forest.
- **Star:** There exist a node such that every link in the network involves i . i is the **center** of the star.

Facts about Trees:

- A connected network is a tree if and only if there has $n - 1$ links.
- In a tree, there is unique path between any two links.

Circles and Complete Networks

- **Complete Network** is one in which all possible links are present, so that $g_{ij} = 1$ for all $i \neq j$.
- **Circle** is a network that has a single cycle and is such that each node has exactly two neighbours.



Neighbourhood

- **Neighbourhood of node i** is a set of nodes that i is linked to.

$$N_i(g) = \{j : g_{ij} = 1\}$$

- Given some nodes S , the **neighbourhood of S** is the union of the neighbourhoods of its members:

$$N_S(g) = \bigcup_{i \in S} N_i(g) = \{j : \exists i \in S, g_{ij} = 1\}$$

- **k -neighbourhood of i** : All nodes that can be reached from i by walks of length no more than k :

$$N_i^k(g) = N_i(g) \cup \left(\bigcup_{j \in N_i(g)} N_j^{k-1}(g) \right)$$

Degree and Network Density

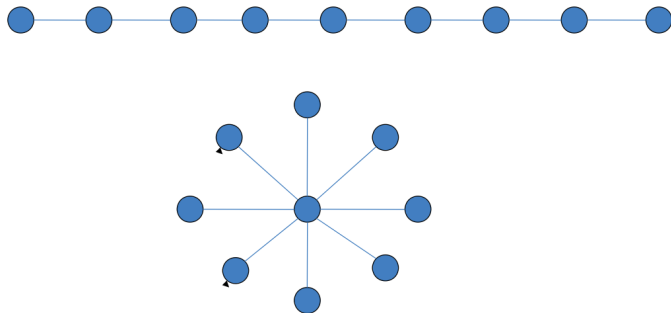
- **Degree**: of a node is the number of links that involves that node:

$$d_j(g) = \#N_i(g) = \#\{j : g_{ij} = 1\}$$

- In terms of **directed networks**:
 - **In-degree**: $d_j(g) = \#\{j : g_{ji} = 1\}$
 - **Out-degree**: $d_j(g) = \#\{j : g_{ij} = 1\}$
- **Network Density**: keeps track of the relative fraction of links present in the network and is equal to **average degree divided by $n - 1$** .

Degree Distribution

- How is degree distributed in the network?
- Average degree only tells us part of the story:



Degree Distribution

- **Degree Distribution** of a network is description of the relative frequency of nodes that have different degrees.
- $P(d)$ is the fraction of nodes that have degree d under distribution p .
 - A **regular network** is one in which all nodes have the same degree.
 - A network is a **regular of degree k** if $P(k) = 1$ and $P(d) = 0 \forall d \neq k$.
- Another example is the **scale-free** or **power** degree distribution:

$$P(d) = cd^{-\gamma}$$

- The relative probabilities of degrees of a fixed relative ratio are the same independent of the scale of those degrees.

$$\frac{P(2)}{P(1)} = \frac{P(20)}{p(10)}$$

Overall Clustering

- What fraction of my friends are friends with each other?
- Overall Clustering:

$$\begin{aligned} Cl(g) &= \frac{\sum_i \#\{jk \in g \mid k \neq j, j \in N_i(g), k \in N_i(g)\}}{\sum_i \#\{jk \mid k \neq j, j \in N_i(g), k \in N_i(g)\}} \\ &= \frac{\sum_i g_{ij}g_{ik}g_{jk}}{\sum_i g_{ij}g_{ik}} \end{aligned}$$

- This is in fact the fraction of fully connected triples out of the potential triples in which at least two links are present.

Individual Clustering

- This measure is computed on node-by-node basis and then averaged across nodes. **Individual Clustering:**

$$\begin{aligned} Cl_i(g) &= \frac{\#\{jk \in g \mid k \neq j, j \in N_i(g), k \in N_i(g)\}}{\#\{jk \mid k \neq j, j \in N_i(g), k \in N_i(g)\}} \\ &= \frac{\sum_{j,k} g_{ij} g_{ik} g_{jk}}{\sum_{j,k} g_{ij} g_{ik}} \end{aligned}$$

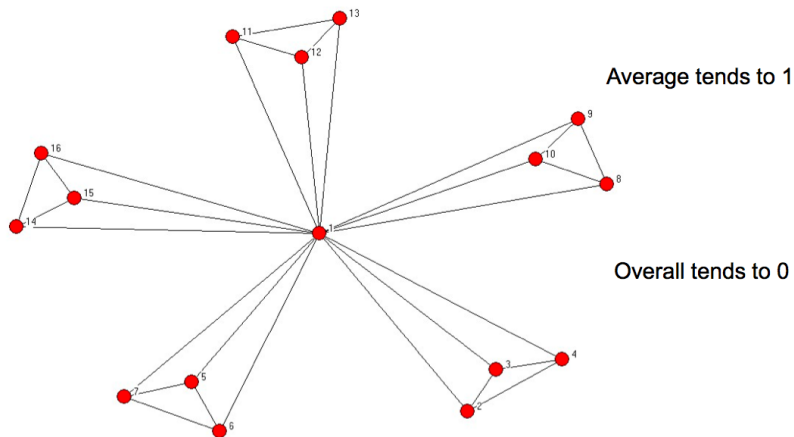
- This looks at all the pairs of nodes that are linked to i and then asks how many of them are connected to each other.
- **Average Clustering:**

$$Cl^{Avg}(g) = \sum_i Cl_i(g) / n$$

Difference In Clustering

- Under **average clustering**, one computes clustering for each node and then averages over all nodes.
- Whereas with **overall clustering**, the average is taken over all triplets.
- Average clustering gives more weight to low-degree nodes than does the clustering coefficient method.
- The average clustering for the Florentine marriage network is $3/20$, where as the overall clustering coefficient is $9/47$.

Difference In Clustering



Position In The Network

How to describe individual characteristics?

- Degree
- Clustering
- Distance to other nodes
- Centrality and Influence

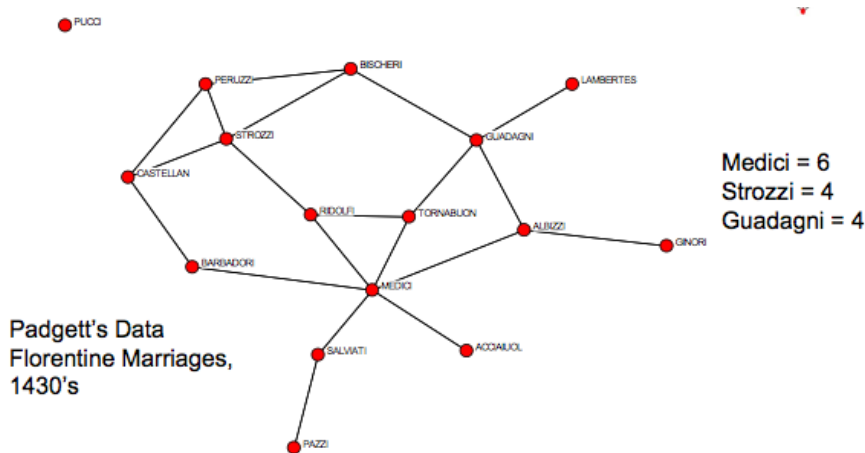
Centrality Measures

- **Degree**: Measure of connectedness.
- **Closeness, Decay**: Ease of reaching other nodes.
- **Betweenness**: Importance as an intermediary, connector.
- **Influence, Prestige, Eigenvectors**: “not what you know, but who you know!”

How “connected” is a node?

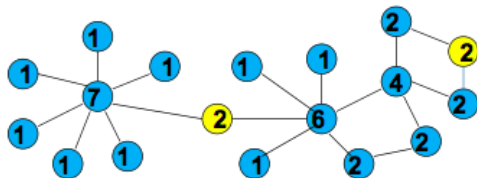
- Degree Centrality captures connectedness.
- Normalize by $n - 1$ Which the most possible number of connections in a network of size n .

Degree Centrality - Examples

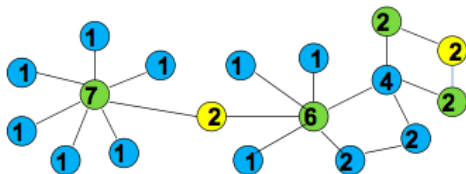


Degree Centrality - Examples

- Failure of degree centrality to capture reach of a node:



- More reach if connected to a 6 and 7 than a 2 and 2?



Eigenvalue Centrality

- **Eigenvalue Centrality** is proportional to the **sum of neighbours' centralities**:

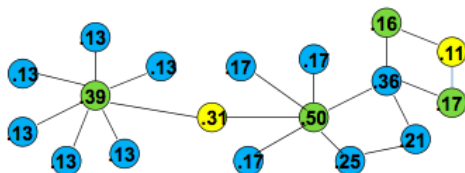
$$C_i \propto \sum_{j \in N_i(g)} C_j$$

- More connections matter, but also accounts for how central they are!

$$\begin{aligned} \alpha C_i &= \sum_{j \in N_i(g)} C_j = \sum_j g_{ij} C_j \\ \Rightarrow \quad \alpha \mathbf{C} &= \mathbf{gC} \end{aligned}$$

- **Google page rank**: score of a page is proportional to the sum of the scores of pages linked to it.

Eigenvalue Centrality - Example



● RUCCI



Padgett's Data
Florentine Marriages,
1430's

Medici = .430
Strozzi = .356
Guadagni = .289
Ridolfi = .341
Tornabuon = .326

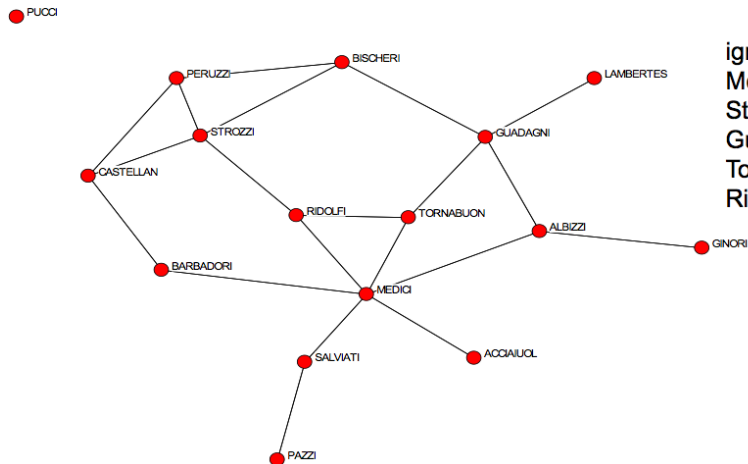
Closeness Centrality

- How close a given node is to any other node?
- **Closeness Centrality**: The inverse of the average distance between i and any other node j :

$$\frac{n - 1}{\sum_{i \neq j} l(i, j)}$$

- $l(i, j)$ is the number of links in the shortest path between i and j .
- Scales directly with distance: twice as far is half as central.

Closeness Centrality



ignoring Pucci:
Medici 14/25
Strozzi 14/32
Guadagni 14/26
Tornabuoni 14/29
Ridolfi 14/28

Decay Centrality

- A richer way of measuring centrality based on closeness is considering a **decay parameter** $0 < \delta < 1$, and consider the proximity between a given node and every other node weighted by this parameter.

- **Decay Centrality:**

$$C_i^d(g) = \sum_{i \neq j} \delta^{l(i,j)}$$

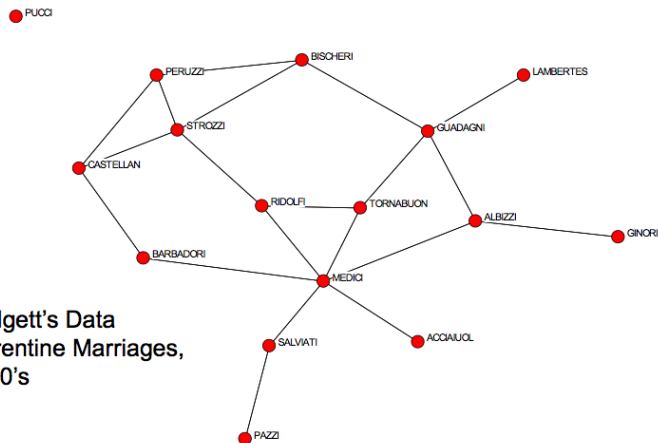
- When $\delta \rightarrow 1$, then $C_i^d(g)$ measures **component size**.
- When $\delta \rightarrow 0$, then the decay centrality gives infinitely more weights to the closer nodes than farther ones, i.e. becomes proportional to **Degree Centrality**.
- For **intermediate values** of δ , closer nodes have higher weights than less close nodes.

Betweenness Centrality

- **Betweenness Centrality** gives a measure of how well situated a node is in terms of the paths that it lies on.
- Let $P(i, j)$ be the number of geodesics between i and j ,
- Let $P_k(i, j)$ be the number of geodesics between i and j that goes through k ,
- Estimate how important k is in terms of connecting i and j , by looking at the ratio:

$$C_i^B(g) = \sum_{i \neq j, k \notin \{i, j\}} \frac{P_k(ij)/P(ij)}{(n-1)(n-2)/2}$$

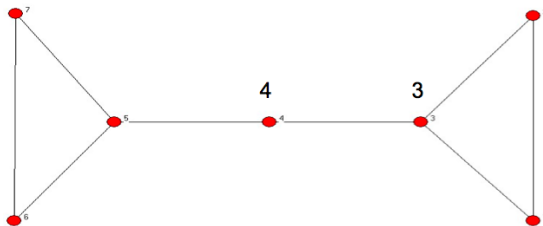
Betweenness Centrality



Medici = .522
Strozzi = .103
Guadagni = .255

Padgett's Data
Florentine Marriages,
1430's

Centrality Measures - Example



	Node 1	Node 3	Node 4
Degree	.33	.50	.33
Closeness	.40	.55	.60
N. Decay $\delta = .5$.50	.67	.67
N. Decay $\delta = .75$.69	.82	.84
N. Decay $\delta = .25$.39	.56	.50
Betweenness	.00	.53	.60
Eigenvector	.47	.63	.54

Centrality Measures

- Degree Centrality: Measures connectedness.
- Closeness and Decay centrality: Measures ease of reaching others.
- Betweenness Centrality: Measures importance as an intermediary and a connector.
- Influence, Prestige, Eigenvectors Centrality: “not what you know, but who you know...”.