

## 7 Some Chromatic Numbers

Let  $G = (V, E)$  be a graph, and let  $\chi(G)$ ,  $\omega(G)$  denote the chromatic, clique numbers of  $G$ .

**Homomorphisms:** If  $G, H$  are graphs a homomorphism from  $G$  to  $H$  is a map  $\phi : V(G) \rightarrow V(H)$  with the property that  $\phi(x)\phi(y)$  is an edge of  $H$  whenever  $xy$  is an edge of  $G$ . If such a map exists, we write  $G \rightarrow H$ . Note that  $G \rightarrow K_n$  if and only if  $\chi(G) \leq K_n$ .

### 7.1 Fractional Colouring

Let  $\mathcal{I}(G)$  denote the collection of all independent sets in  $G$  and let  $A$  be the  $V \times \mathcal{I}(G)$  incidence matrix (i.e.  $A_{v,S}$  is 1 if  $v$  is in the independent set  $S$  and 0 otherwise). Now we can express  $\chi(G)$  and  $\omega(G)$  with the following integer programs.

$$\begin{aligned}\chi(G) &= \min\{1^\top x : x \in \mathbb{Z}_+^{\mathcal{I}(G)} \text{ and } Ax \geq 1\} \\ \omega(G) &= \min\{y^\top 1 : y \in \mathbb{Z}_+^V \text{ and } y^\top A \leq 1\}\end{aligned}$$

Relaxing these integrality constraints yield the *fractional clique number*  $\omega_f$ , and *fractional chromatic number*  $\chi_f$ , expressed as linear programs.

$$\begin{aligned}\chi_f(G) &= \min\{1^\top x : x \in \mathbb{R}_+^{\mathcal{I}(G)} \text{ and } Ax \geq 1\} \\ \omega_f(G) &= \min\{y^\top 1 : y \in \mathbb{R}_+^V \text{ and } y^\top A \leq 1\}\end{aligned}$$

Now, by LP duality we have the following string of inequalities

$$\omega(G) \leq \omega_f(G) = \chi_f(G) \leq \chi(G).$$

**Note:** Although linear programs can be solved in polynomial time, the size of the matrix  $A$  is generally exponential in the size of  $G$ , so this does not allow us to compute  $\chi_f$  efficiently. In fact, it is NP-hard to compute  $\chi_f(G)$ .

**Kneser Graphs:** Let  $\binom{[n]}{k}$  denote the set of all  $k$  element subsets of  $\{1, \dots, n\}$ . The Kneser graph  $Kn(n, k)$  has vertex set  $\binom{[n]}{k}$ , with two vertices adjacent if they are disjoint.

**Theorem 7.1**  $\chi_f(G) \leq \frac{p}{q}$  if and only if  $G \rightarrow Kn(\ell p, \ell q)$  for some  $\ell \in \mathbb{Z}$ .

*Proof:* "only if": If  $\chi_f(G) \leq \frac{p}{q}$  then there is a rational vector  $w \in \mathbb{Q}^{\mathcal{I}(G)}$  so that  $\sum\{w_I : v \in I\} \geq 1$  for every  $v \in V$  and  $w^\top \mathbf{1} = \frac{p}{q}$ . Choose an integer  $\ell$  so that  $(\ell q)w \in \mathbb{Z}^{\mathcal{I}(G)}$ . Now, form a sequence  $L$  of independent sets so that  $(\ell q)w_I$  is the number of times  $I$  appears in  $L$ . By construction, the length of  $L$  is equal to the sum of the entries in  $(\ell q)w$  which is precisely  $\ell p$ . Furthermore, for every  $v \in V$  we have  $\sum\{(\ell q)w_I : v \in I\} \geq \ell q$ , so  $v$  appears in  $\geq \ell q$  of the terms of  $L$ . By possibly removing  $v$  from some terms of  $L$ , we may arrange that every vertex appears in exactly  $\ell q$  terms. Now, let  $\phi : V(G) \rightarrow \binom{[\ell p]}{\ell q}$  be given by the rule

$$\phi(v) = \{j \in \{1, \dots, \ell p\} : v \in L_j\}$$

It follows that  $\phi$  is a homomorphism from  $G$  to  $Kn(\ell p, \ell q)$  as required.

"if": To prove this it suffices to show that  $\chi_f(Kn(p, q)) \leq \frac{p}{q}$ , since the composition of a homomorphism and a fractional  $\frac{p}{q}$ -colouring is another fractional  $\frac{p}{q}$ -colouring. A fractional  $\frac{p}{q}$ -colouring of  $Kn(p, q)$  is given by assigning weight  $\frac{1}{q}$  to each independent set of the form  $T_i = \{S \in \binom{[p]}{q} : i \in S\}$ .  $\square$

## 7.2 Circular Colouring

For  $t \in \mathbb{R}$ , a *circular  $t$ -colouring* of  $G$  is a map  $\phi : V \rightarrow S^1$  so that the angle between  $\phi(x)$  and  $\phi(y)$  is  $\geq \frac{2\pi}{t}$  whenever  $x$  and  $y$  are adjacent. The *circular chromatic number* of  $G$  is

$$\chi_c(G) = \inf\{t \in \mathbb{R} : G \text{ has a circular } t\text{-colouring}\}.$$

$\mathbf{K}_{n/k}$ : If  $n, k$  are positive integers, we let  $K_{n/k}$  denote a graph consisting of  $n$  vertices in a cyclic order, with two vertices adjacent if they have distance  $\geq k$  in this ordering.

### Theorem 7.2

- (i)  $\lceil \chi_c(G) \rceil = \chi(G)$
- (ii)  $\chi_c(G) \leq \frac{n}{k}$  if and only if  $G \rightarrow K_{n/k}$ .

### 7.3 Vector Colouring

For  $t \in \mathbb{R}$ , a *vector  $t$ -colouring* of  $G$  is a mapping  $x : V \rightarrow S^m$  with the property that  $x(i) \cdot x(j) \leq -\frac{1}{t-1}$  whenever  $ij \in E$ . The *vector chromatic number* of  $G$  is

$$\chi_v(G) = \inf\{t \in \mathbb{R} : G \text{ has a vector } t\text{-colouring}\}.$$

**Note:**  $\chi_v(G)$  can be computed (efficiently!) with the following SDP (here  $X \in \mathbb{R}^{V \times V}$ )

$$\begin{aligned} & \min s \\ & X \succeq 0 \\ & X_{ii} = 1 \text{ for every } i \in V \\ & X_{ij} \leq s \text{ whenever } ij \in E \end{aligned}$$

**Theorem 7.3**  $\omega(G) \leq \chi_v(G) \leq \chi_f(G) \leq \chi_c(G) \leq \chi(G)$

*Proof:*  $\omega(G) \leq \chi_v(G)$ : Let  $x : V \rightarrow S^n$  be a vector  $t$ -colouring of  $G$  and choose a clique  $S \subseteq V$  with  $|S| = \omega(G)$ . Now we have

$$0 \leq \left( \sum_{i \in S} x(i) \right) \cdot \left( \sum_{i \in S} x(i) \right) = |S| + 2 \sum_{i,j \in S: i \neq j} x(i) \cdot x(j) \leq |S| + |S|(|S| - 1) \left( -\frac{1}{t-1} \right)$$

from which we deduce  $t \geq |S| = \omega(G)$  as desired.

$\chi_v(G) \leq \chi_f(G)$ : It suffices to show that  $\chi_v(Kn(n, k)) \leq \frac{n}{k}$  as the composition of a homomorphism from  $G$  to  $Kn(n, k)$  with a vector  $t$ -colouring of  $Kn(n, k)$  is a vector  $t$ -colouring of  $G$ . For every vertex  $S \subseteq \{1, 2, \dots, n\}$  in we define  $w(S) \in \mathbb{R}^n$  as follows:

$$w(S)_i = \begin{cases} k - n & \text{if } i \in S \\ k & \text{if } i \notin S \end{cases}$$

Now  $w$  assigns each vertex  $S$  a vector  $w(S)$  with  $\|w(S)\|^2 = k(n-k)^2 + (n-k)k^2 = k(n-k)n$  and if  $S, T$  are adjacent vertices, then  $w(S) \cdot w(T) = (2k)k(k-n) + (n-2k)(k^2) = -k^2n$  so now setting  $x = \frac{1}{\sqrt{k(n-k)n}}w$  we have that  $\|x(S)\| = 1$  for every  $S$  and  $x(S) \cdot x(T) = \frac{-k}{n-k} = \frac{-1}{n/k-1}$  whenever  $S$  and  $T$  are adjacent, so  $x$  is a vector  $\frac{n}{k}$  colouring.

$\chi_f(G) \leq \chi_c(G)$ : It suffices to show that  $\chi_f(K_{n/k}) \leq \frac{n}{k}$  since the composition of a homomorphism from  $G$  to  $K_{n/k}$  and a fractional  $\frac{n}{k}$ -colouring of  $K_{n/k}$  is a fractional  $\frac{n}{k}$ -colouring of  $G$ . A fractional  $\frac{n}{k}$ -colouring of  $K_{n/k}$  is given by assigning weight  $\frac{1}{k}$  to each of the  $n$  intervals of length  $k$  in the cyclic ordering of the vertices.

$\chi_c(G) \leq \chi(G)$ : This follows immediately from Theorem 7.2.  $\square$